Session Agenda

- Session Overview
- ER and EER to Relational Mapping
- Database Design Methodology and UML
- Mapping Relational Design to ER/EER Case Study
- Summary & Conclusion
What is the class about?

- Course description and syllabus:
  - [http://www.nyu.edu/classes/jcf/CSCI-GA.2433-001](http://www.nyu.edu/classes/jcf/CSCI-GA.2433-001)
  - [http://cs.nyu.edu/courses/spring16/CSCI-GA.2433-001/](http://cs.nyu.edu/courses/spring16/CSCI-GA.2433-001/)

- Textbooks:
    - Ramez Elmasri and Shamkant Navathe
    - Pearson
Icons / Metaphors

- Information
- Common Realization
- Knowledge/Competency Pattern
- Governance
- Alignment
- Solution Approach
1. Session Overview
2. ER and EER to Relational Mapping
3. Database Design Methodology and UML
4. Mapping Relational Design to ER/EER Case Study
5. Summary and Conclusion
Agenda

- Sets
- Relations and tables
- Relational schema
- Primary keys
- Relational Database Design Using ER-to-Relational Mapping
- Mapping EER Model Constructs to Relations
- Design a relational database schema
  - Based on a conceptual schema design
- Seven-step algorithm to convert the basic ER model constructs into relations
- Additional steps for EER model
In this unit, we learn the semantics of specifying a relational database, later we will learn the syntax of SQL for doing this.

The basic “datatype”, or “variable” of a relational database is a relation.

In this unit, such a variable will be a set.

Later, we will extend this, and such a variable will be a multiset.

In SQL, such a variable is called a table.

We may use the term table for a relation in this unit too.
- We will not use axiomatic set theory
- A **set** is a “bag” of elements, some/all of which could be sets themselves and a binary relationship “is element of” denoted by $\in$, such as $2 \in \{2, 5, 3, 7\}$, $\{2,8\} \in \{2, \{2, 8\}, 5, 3, 7\}$,
- You cannot specify
  - How many times an element appears in a set (if you could, this would be a **multiset**)
  - In which position an element appears (if you could, this would be a **sequence**)
- Therefore, as sets: $\{2, 5, 3, 7\} = \{2, 7, 5, 3, 5, 3, 3\}$
- Note: in many places you will read: “an element can appear in a set only once”
This is not quite right. And it is important not to assume this, as we will see in the next unit
Two sets $A$ and $B$ are equal iff (if and only if) they have the same elements.

In other words, for every $x$: $x$ is an element of $A$ iff (if and only if) $x$ is an element of $B$.

“More mathematically,”

$\forall x \{ x \in A \iff x \in B \}$ if and only if $A = B$.

Therefore, as sets: $\{2, 5, 3, 7\} = \{2, 7, 5, 3, 5, 3, 3\}$.

This reiterates what we have said previously.
Consider a table, with a fixed number of columns where elements of each column are drawn from some specific domain.

The columns are labeled and the labels are distinct.

We will consider such a table to be *a set of rows* (another word for “row”: *tuple*).

Here is an example of a table $S$ of two columns $A$ and $B$:

<table>
<thead>
<tr>
<th></th>
<th>A</th>
<th>B</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>a</td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>b</td>
<td>3</td>
<td></td>
</tr>
<tr>
<td>c</td>
<td>4</td>
<td></td>
</tr>
<tr>
<td>d</td>
<td>3</td>
<td></td>
</tr>
</tbody>
</table>

A *relation* is such a table.

We will also write $S(A,B)$ for table $S$ with columns $A$ and $B$. 

What we saw was an *instance* (current value for a relation with the defined columns and domains)

To specify this relation in general (not the specific instance) we need to talk about a *relational schema*

A relational schema defines a constant number of relations, one or more
Relational Schema

- Here is an informal, but complete, description of what is a relational schema of one relation
- We want to define a structure for some table
  1. We give it a name (we had S)
  2. We chose the number of columns (we had 2) and give them distinct names (we had A and B)
  3. We decide on the domains of elements in the columns (we had letters for A and integers for B)
  4. We decide on constraints, if any, on the permitted values (we had that any two rows that are equal on A must be equal on B)
Let’s verify

- $A$ all lower case letters in English
- $B$ all positive integers less than 100
- $S(A,B)$ satisfies the condition that any two tuples that are equal on $A$ must also be equal on $B$

Our example was an instance of this relational schema

<table>
<thead>
<tr>
<th>S</th>
<th>A</th>
<th>B</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>a</td>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>b</td>
<td>b</td>
<td>3</td>
</tr>
<tr>
<td>c</td>
<td>c</td>
<td>4</td>
</tr>
<tr>
<td>d</td>
<td>d</td>
<td>3</td>
</tr>
</tbody>
</table>
Since relations are *sets* of tuples, *the following two relations are equal* (are really one relation written in two different ways)

(This is a different example, not an instance of the previous relational schema)
Since *the positions* in the tuple (1st, 2nd, etc.) are labeled with the column headings, *the following two relations are equal* (are really one relation written in two different ways)

<table>
<thead>
<tr>
<th>S</th>
<th>A</th>
<th>B</th>
</tr>
</thead>
<tbody>
<tr>
<td>a</td>
<td>2</td>
<td>a</td>
</tr>
<tr>
<td>a</td>
<td>56</td>
<td>a</td>
</tr>
<tr>
<td>b</td>
<td>2</td>
<td>a</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>S</th>
<th>B</th>
<th>A</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>56</td>
<td>a</td>
</tr>
<tr>
<td></td>
<td>2</td>
<td>a</td>
</tr>
<tr>
<td>2</td>
<td>56</td>
<td>a</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>b</td>
</tr>
<tr>
<td>56</td>
<td>a</td>
<td>a</td>
</tr>
<tr>
<td>2</td>
<td>a</td>
<td>a</td>
</tr>
<tr>
<td>56</td>
<td>a</td>
<td>a</td>
</tr>
</tbody>
</table>
To specify relations, it is enough to do what we have done above.

As long as we understand what are the domains for the columns, the following are formally fully specified relations:

- Relational (schema) $P(\text{Name}, \text{SSN}, \text{DOB}, \text{Grade})$ with some (not specified, but we should have done it) domains for attributes
- Relational (schema) $Q(\text{Grade}, \text{Salary})$ with some (not specified, but we should have done it) domains for attributes

<table>
<thead>
<tr>
<th>P</th>
<th>Name</th>
<th>SSN</th>
<th>DOB</th>
<th>Grade</th>
</tr>
</thead>
<tbody>
<tr>
<td>A</td>
<td>121</td>
<td>2367</td>
<td>2</td>
<td></td>
</tr>
<tr>
<td>B</td>
<td>101</td>
<td>3498</td>
<td>4</td>
<td></td>
</tr>
<tr>
<td>C</td>
<td>106</td>
<td>2987</td>
<td>2</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Q</th>
<th>Grade</th>
<th>Salary</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>90</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>80</td>
<td></td>
</tr>
<tr>
<td>3</td>
<td>70</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>70</td>
<td></td>
</tr>
</tbody>
</table>
But we will do more. We will specify, as appropriate for the schema:

- **Primary keys**
- **Keys** (beyond primary)
- **Foreign keys** and what they reference (we will see soon what this means)
- Additional constraints

Some of the constraints involve more than one relation

The above most important structurally

Later, when we talk about SQL DDL, we will specify additional properties
Consider relation (schema) Person(FN, LN, Grade, YOB)

Instance:

<table>
<thead>
<tr>
<th>Person</th>
<th>FN</th>
<th>LN</th>
<th>Grade</th>
<th>YOB</th>
</tr>
</thead>
<tbody>
<tr>
<td>John</td>
<td>Smith</td>
<td>8</td>
<td>1976</td>
<td></td>
</tr>
<tr>
<td>Lakshmi</td>
<td>Smith</td>
<td>9</td>
<td>1981</td>
<td></td>
</tr>
<tr>
<td>John</td>
<td>Smith</td>
<td>8</td>
<td>1976</td>
<td></td>
</tr>
<tr>
<td>John</td>
<td>Yao</td>
<td>9</td>
<td>1992</td>
<td></td>
</tr>
</tbody>
</table>

- We are told that any two tuples that are equal on both FN and LN are (completely) equal
  - We have some tuples appearing multiple times: this is just for clarifying that this permitted in the definition, we do not discuss here why we would have the same tuple more than one time (we will talk about this later)

- This is a property of every possible instance of Person in our application—we are told this

- Then (FN, LN) is a superkey of Person, and in fact a key, because neither FN nor LN by themselves are sufficient (we are told that too)
Consider relation (schema) Q(Grade, Salary)

Example:

<table>
<thead>
<tr>
<th>Pay</th>
<th>Grade</th>
<th>Salary</th>
</tr>
</thead>
<tbody>
<tr>
<td>8</td>
<td></td>
<td>128</td>
</tr>
<tr>
<td>9</td>
<td></td>
<td>139</td>
</tr>
<tr>
<td>7</td>
<td></td>
<td>147</td>
</tr>
</tbody>
</table>

We are told that for any instance of Pay, any two tuples that are equal on Grade are (completely) equal.

Of course, if each Grade appears in only one tuple, this is automatically true.

Then, similarly to before, Grade is a key.

What about Salary, is this a key also?

No, because we are not told that any two tuples that are equal on Salary are equal on Grade in every instance of Pay.
A set of columns in a relation is a **superkey** if and only any two tuples that are equal on the elements of these columns are (completely equal)

- **A relation always has at least one superkey**
- The set of all the attributes is a superkey
- Because any two tuples that are equal on all attributes are completely equal
- A minimal superkey, is a key
- A relation always has at least one key (start with any superkey and remove unnecessary columns)
- There may be more than one key
- Exactly one key is chosen as **primary key**
- Other keys are just keys
- Sometimes they are called **candidate keys** (as they are candidates for the primary key, though not chosen)
- We will underline the attributes of the chosen primary key.
- Returning to the City example: City(Longitude, Latitude, Country, State, Name, Size)

We can have
- City(Longitude, Latitude, Country, State, Name, Size)
  - This implies that Longitude, Latitude form a primary key
  - We also have a candidate key: Country, State, Name

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- City(Longitude, Latitude, Country, State, Name, Size)
  - This implies that Country, State, Name form a primary key
  - We also have a candidate key: Longitude, Latitude
A relational database is basically a set of relations and is an instance of a relational schema.
ER conceptual schema diagram for the COMPANY database
GOALS during Mapping

- Preserve all information (that includes all attributes)
- Maintain the constraints to the extent possible (Relational Model cannot preserve all constraints- e.g., max cardinality ratio such as 1:10 in ER; exhaustive classification into subtypes, e.g., STUDENTS are specialized into Domestic and Foreign)
- Minimize null values

*The mapping procedure described has been implemented in many commercial tools.*
- **ER-to-Relational Mapping Algorithm**
  - Step 1: Mapping of Regular Entity Types
  - Step 2: Mapping of Weak Entity Types
  - Step 3: Mapping of Binary 1:1 Relation Types
  - Step 4: Mapping of Binary 1:N Relationship Types
  - Step 5: Mapping of Binary M:N Relationship Types
  - Step 6: Mapping of Multivalued attributes.
  - Step 7: Mapping of N-ary Relationship Types.

- **Mapping EER Model Constructs to Relations**
  - Step 8: Options for Mapping Specialization or Generalization.
  - Step 9: Mapping of Union Types (Categories).
COMPANY database example
- Assume that the mapping will create tables with simple single-valued attributes

Step 1: Mapping of Regular Entity Types.
- For each regular (strong) entity type E in the ER schema, create a relation R that includes all the simple attributes of E.
- Choose one of the key attributes of E as the primary key for R.
- If the chosen key of E is composite, the set of simple attributes that form it will together form the primary key of R.

Example: We create the relations EMPLOYEE, DEPARTMENT, and PROJECT in the relational schema corresponding to the regular entities in the ER diagram.
- SSN, DNUMBER, and PNUMBER are the primary keys for the relations EMPLOYEE, DEPARTMENT, and PROJECT as shown.
Step 2: Mapping of Weak Entity Types

- For each weak entity type W in the ER schema with owner entity type E, create a relation R & include all simple attributes (or simple components of composite attributes) of W as attributes of R.
- Also, include as foreign key attributes of R the primary key attribute(s) of the relation(s) that correspond to the owner entity type(s).
- The primary key of R is the \textit{combination of} the primary key(s) of the owner(s) and the partial key of the weak entity type W, if any.

Example: Create the relation DEPENDENT in this step to correspond to the weak entity type DEPENDENT.

- Include the primary key SSN of the EMPLOYEE relation as a foreign key attribute of DEPENDENT (renamed to ESSN).
- The primary key of the DEPENDENT relation is the combination \{ESSN, DEPENDENT\_NAME\} because DEPENDENT\_NAME is the partial key of DEPENDENT.
**Figure 9.3**
Illustration of some mapping steps.
b. Additional *weak entity* relation after step 2.

(a) **EMPLOYEE**

<table>
<thead>
<tr>
<th>Fname</th>
<th>Minit</th>
<th>Lname</th>
<th>Ssn</th>
<th>Bdate</th>
<th>Address</th>
<th>Sex</th>
<th>Salary</th>
</tr>
</thead>
</table>

(b) **DEPARTMENT**

<table>
<thead>
<tr>
<th>Dname</th>
<th>Dnumber</th>
</tr>
</thead>
</table>

(c) **PROJECT**

<table>
<thead>
<tr>
<th>Pname</th>
<th>Pnumber</th>
<th>Plocation</th>
</tr>
</thead>
</table>

(b) **DEPENDENT**

<table>
<thead>
<tr>
<th>Essn</th>
<th>Dependent_name</th>
<th>Sex</th>
<th>Bdate</th>
<th>Relationship</th>
</tr>
</thead>
</table>

(c) **WORKS_ON**

<table>
<thead>
<tr>
<th>Essn</th>
<th>Pno</th>
<th>Hours</th>
</tr>
</thead>
</table>

(d) **DEPT_LOCATION**

<table>
<thead>
<tr>
<th>Dnumber</th>
<th>Dlocation</th>
</tr>
</thead>
</table>
Step 3: Mapping of Binary 1:1 Relation Types

- For each binary 1:1 relationship type R in the ER schema, identify the relations S and T that correspond to the entity types participating in R.

There are three possible approaches:

1. **Foreign Key (2 relations) approach:** Choose one of the relations—say S—and include a foreign key in S the primary key of T. It is better to choose an entity type with total participation in R in the role of S.
   - Example: 1:1 relation MANAGES is mapped by choosing the participating entity type DEPARTMENT to serve in the role of S, because its participation in the MANAGES relationship type is total.

2. **Merged relation (1 relation) option:** An alternate mapping of a 1:1 relationship type is possible by merging the two entity types and the relationship into a single relation. This may be appropriate when both participations are total.

3. **Cross-reference or relationship relation (3 relations) option:** The third alternative is to set up a third relation R for the purpose of cross-referencing the primary keys of the two relations S and T representing the entity types.
Step 4: Mapping of Binary 1:N Relationship Types.

For each regular binary 1:N relationship type $R$, identify the relation $S$ that represent the participating entity type at the $N$-side of the relationship type.

Include as foreign key in $S$ the primary key of the relation $T$ that represents the other entity type participating in $R$.

Include any simple attributes of the 1:N relation type as attributes of $S$.

Example: 1:N relationship types WORKS_FOR, CONTROLS, and SUPERVISION in the figure.

For WORKS_FOR we include the primary key DNUMBER of the DEPARTMENT relation as foreign key in the EMPLOYEE relation and call it DNO.

An alternative approach is to use a Relationship relation (cross referencing relation) as in the third option for binary 1:1 relationships – this is rarely done.
Step 5: Mapping of Binary M:N Relationship Types.

» For each regular binary M:N relationship type R, create a new relation S to represent R. This is a relationship relation.

» Include as foreign key attributes in S the primary keys of the relations that represent the participating entity types; their combination will form the primary key of S.

» Also include any simple attributes of the M:N relationship type (or simple components of composite attributes) as attributes of S.

Example: The M:N relationship type WORKS_ON from the ER diagram is mapped by creating a relation WORKS_ON in the relational database schema.

» The primary keys of the PROJECT and EMPLOYEE relations are included as foreign keys in WORKS_ON and renamed PNO and ESSN, respectively.

» Attribute HOURS in WORKS_ON represents the HOURS attribute of the relation type. The primary key of the WORKS_ON relation is the combination of the foreign key attributes \{ESSN, PNO\}.
Step 6: Mapping of Multivalued attributes.

- For each multivalued attribute A, create a new relation R.
- This relation R will include an attribute corresponding to A, plus the primary key attribute K as a foreign key in R of the relation that represents the entity type of relationship type that has A as an attribute.
- The primary key of R is the combination of A and K. If the multivalued attribute is composite, we include its simple components.

Example: The relation DEPT_LOCATIONS is created.

- The attribute DLOCATION represents the multivalued attribute LOCATIONS of DEPARTMENT, while DNUMBER-as foreign key-represents the primary key of the DEPARTMENT relation.
- The primary key of R is the combination of {DNUMBER, DLOCATION}. 
Step 7: Mapping of N-ary Relationship Types.

- For each n-ary relationship type R, where n>2, create a new relationship S to represent R.
- Include as foreign key attributes in S the primary keys of the relations that represent the participating entity types.
- Also include any simple attributes of the n-ary relationship type (or simple components of composite attributes) as attributes of S.

Example: The relationship type SUPPY in the ER on the next slide.

- This can be mapped to the relation SUPPLY shown in the relational schema, whose primary key is the combination of the three foreign keys {SNAME, PARTNO, PROJNAME}.
Result of Mapping the COMPANY ER Schema into a RDB Schema
Ternary Relationship: SUPPLY

(a) SUPPLIER → SUPPLY → PROJECT

- SName
- Quantity
- ProjName

- PartNo
- PART
Mapping the n-ary Relationship Type SUPPLY

SUPPLIER

Sname  ...  

PROJECT

Proj_name  ...  

PART

Part_no  ...  

SUPPLY

Sname  Proj_name  Part_no  Quantity
### Table 9.1  Correspondence between ER and Relational Models

<table>
<thead>
<tr>
<th>ER MODEL</th>
<th>RELATIONAL MODEL</th>
</tr>
</thead>
<tbody>
<tr>
<td>Entity type</td>
<td><em>Entity</em> relation</td>
</tr>
<tr>
<td>1:1 or 1:N relationship type</td>
<td>Foreign key (or <em>relationship</em> relation)</td>
</tr>
<tr>
<td>M:N relationship type</td>
<td><em>Relationship</em> relation and <em>two</em> foreign keys</td>
</tr>
<tr>
<td><em>n</em>-ary relationship type</td>
<td><em>Relationship</em> relation and <em>n</em> foreign keys</td>
</tr>
<tr>
<td>Simple attribute</td>
<td><em>Attribute</em></td>
</tr>
<tr>
<td>Composite attribute</td>
<td>Set of simple component attributes</td>
</tr>
<tr>
<td>Multivalued attribute</td>
<td><em>Relation</em> and foreign key</td>
</tr>
<tr>
<td>Value set</td>
<td><em>Domain</em></td>
</tr>
<tr>
<td>Key attribute</td>
<td>Primary (or secondary) key</td>
</tr>
</tbody>
</table>
- In a relational schema relationship, types are not represented explicitly
  - Represented by having two attributes $A$ and $B$: one a primary key and the other a foreign key
Step 8: Options for Mapping Specialization or Generalization

Convert each specialization with m subclasses \{S_1, S_2, \ldots, S_m\} and generalized superclass \(C\), where the attributes of \(C\) are \{k, a_1, \ldots, a_n\} and \(k\) is the (primary) key, into relational schemas using one of the four following options:

- Option 8A: Multiple relations - Superclass and subclasses
- Option 8B: Multiple relations - Subclass relations only
- Option 8C: Single relation with one type attribute
- Option 8D: Single relation with multiple type attributes
Option 8A: Multiple relations-Superclass and subclasses

Create a relation L for C with attributes \( \text{Attributes}(L) = \{k,a_1,\ldots,a_n\} \) and \( \text{PK}(L) = k \). Create a relation \( L_i \) for each subclass \( S_i \), \( 1 < i < m \), with the attributes \( \text{Attributes}(L_i) = \{k\} \cup \{\text{attributes of } S_i\} \) and \( \text{PK}(L_i) = k \). This option works for any specialization (total or partial, disjoint or overlapping).

Option 8B: Multiple relations-Subclass relations only

Create a relation \( L_i \) for each subclass \( S_i \), \( 1 < i < m \), with the attributes \( \text{Attributes}(L_i) = \{\text{attributes of } S_i\} \cup \{k,a_1,\ldots,a_n\} \) and \( \text{PK}(L_i) = k \). This option only works for a specialization whose subclasses are total (every entity in the superclass must belong to (at least) one of the subclasses).
Option 8C: Single relation with one type attribute

Create a single relation L with attributes $\text{Attrs}(L) = \{k, a_1, \ldots a_n\} \cup \{\text{attributes of } S_1\} \cup \ldots \cup \{\text{attributes of } S_m\} \cup \{t\}$ and $\text{PK}(L) = k$. The attribute t is called a type (or **discriminating**) attribute that indicates the subclass to which each tuple belongs.

Option 8D: Single relation with multiple type attributes

Create a single relation schema L with attributes $\text{Attrs}(L) = \{k, a_1, \ldots a_n\} \cup \{\text{attributes of } S_1\} \cup \ldots \cup \{\text{attributes of } S_m\} \cup \{t_1, t_2, \ldots, t_m\}$ and $\text{PK}(L) = k$. Each $t_i$, $1 < i < m$, is a Boolean type attribute indicating whether a tuple belongs to the subclass $S_i$. 
Mapping of Shared Subclasses (Multiple Inheritance)

- Apply any of the options discussed in step 8 to a shared subclass

![Diagram of mapping EER specialization lattice using multiple options]

**Figure 9.6**
Mapping the EER specialization lattice in Figure 8.8 using multiple options.
See next slide options for mapping specialization or generalization

(a) Mapping the EER schema in Figure 4.4 using option 8A
(b) Mapping the EER schema in Figure 4.3(b) using option 8B
(c) Mapping the EER schema in Figure 4.4 using option 8C
(d) Mapping Figure 4.5 using option 8D with Boolean type fields Mflag and Pflag
Figure 4.4

EER Diagram Notation for Attribute-Defined Specialization on JobType
Mapping the EER schema in Figure 4.4 using option 8A

(a) EMPLOYEE

<table>
<thead>
<tr>
<th>SSN</th>
<th>FName</th>
<th>MInit</th>
<th>LName</th>
<th>BirthDate</th>
<th>Address</th>
<th>JobType</th>
</tr>
</thead>
</table>

SECRETARY

| SSN | TypingSpeed |

TECHNICIAN

| SSN | TGrade |

ENGINEER

| SSN | EngType |
(c) **EMPLOYEE**

<table>
<thead>
<tr>
<th>SSN</th>
<th>FName</th>
<th>MInit</th>
<th>LName</th>
<th>BirthDate</th>
<th>Address</th>
<th>JobType</th>
<th>TypingSpeed</th>
<th>TGrade</th>
</tr>
</thead>
</table>

Mapping the EER schema in Figure 4.4 using option 8C
Generalizing CAR and TRUCK into the superclass VEHICLE

Figure 4.3 (b)
Mapping the EER schema in Figure 4.3 (b) using option 8B

<table>
<thead>
<tr>
<th>CAR</th>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>VehicleId</td>
<td>LicensePlateNo</td>
<td>Price</td>
<td>MaxSpeed</td>
<td>NoOfPassengers</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>TRUCK</th>
<th></th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
<tbody>
<tr>
<td>VehicleId</td>
<td>LicensePlateNo</td>
<td>Price</td>
<td>NoOfAxles</td>
<td></td>
</tr>
</tbody>
</table>


An overlapping (non-disjoint) specialization.

Figure 4.5
(d) **PART**

<table>
<thead>
<tr>
<th>PartNo</th>
<th>Description</th>
<th>MFlag</th>
<th>DrawingNo</th>
<th>ManufactureDate</th>
<th>BatchNo</th>
<th>PFlag</th>
<th>SupplierName</th>
<th>ListPrice</th>
</tr>
</thead>
</table>

Mapping Fig. 4.5 via option 8B with Boolean Type Fields Mflag & Pflag
### Figure 9.5

#### Different Options for Mapping Generalization Hierarchies - Summary

### EMPLOYEE

<table>
<thead>
<tr>
<th>Ssn</th>
<th>Fname</th>
<th>Minit</th>
<th>Lname</th>
<th>Birth_date</th>
<th>Address</th>
<th>Job_type</th>
</tr>
</thead>
</table>

- **SECRETARY**
  - | Ssn | Typing_speed |

- **TECHNICIAN**
  - | Ssn | Tgrade |

- **ENGINEER**
  - | Ssn | Eng_type |

### CAR

| Vehicle_id | License_plate_no | Price | Max_speed | No_of_passengers |

### TRUCK

| Vehicle_id | License_plate_no | Price | No_of_axles | Tonnage |

### EMPLOYEE

| Ssn | Fname | Minit | Lname | Birth_date | Address | Job_type | Typing_speed | Tgrade | Eng_type |

### PART

| Part_no | Description | Mflag | Drawing_no | Manufacture_date | Batch_no | Pflag | Supplier_name | List_price |

---

**Figure 9.5**
Mapping of Shared Subclasses (Multiple Inheritance)

» A shared subclass, such as STUDENT_ASSISTANT, is a subclass of several classes, indicating multiple inheritance. These classes must all have the same key attribute; otherwise, the shared subclass would be modeled as a category.

» We can apply any of the options discussed in Step 8 to a shared subclass, subject to the restriction discussed in Step 8 of the mapping algorithm. Below both 8C and 8D are used for the shared class STUDENT_ASSISTANT.
Specialization Lattice with Multiple Inheritance for UNIVERSITY DB

Figure 4.7
Mapping EER Specialization lattice in Figure 4.7 Using Multiple Options

PERSON

| Ssn | Name | Birth_date | Sex | Address |

EMPLOYEE

| Ssn | Salary | Employee_type | Position | Rank | Percent_time | Ra_flag | Ta_flag | Project | Course |

ALUMNUS

| Ssn |

ALUMNUS_DEGREES

| Ssn | Year | Degree | Major |

STUDENT

| Ssn | Major_dept | Grad_flag | Undergrad_flag | Degree_program | Class | Student_assist_flag |
Step 9: Mapping of Union Types (Categories).

» For mapping a category whose defining superclass have different keys, it is customary to specify a new key attribute, called a surrogate key, when creating a relation to correspond to the category.

» In the example below we can create a relation OWNER to correspond to the OWNER category and include any attributes of the category in this relation. The primary key of the OWNER relation is the surrogate key, which we called OwnerId.
Sample Mapping of EER Categories to Relations

**Figure 9.7**
Mapping the EER categories (union types) in Figure 8.8 to relations.
Two categories (union types): OWNER and REGISTERED_VEHICLE

Figure 4.8
Mapping the EER Categories (Union Types) in Figure 4.8 to Relations

PERSON
- Ssn
- Driver_license_no
- Name
- Address
- Owner_id

BANK
- Bname
- Baddress
- Owner_id

COMPANY
- Cname
- Caddress
- Owner_id

OWNER
- Owner_id

REGISTERED_VEHICLE
- Vehicle_id
- License_plate_number

CAR
- Vehicle_id
- Cstyle
- Cmake
- Cmodel
- Cyear

TRUCK
- Vehicle_id
- Tmake
- Tmodel
- Tonnage
- Tyear

OWNS
- Owner_id
- Vehicle_id
- Purchase_date
- Lien_or_regular

Figure 9.7
Exercise 9.4: Map this schema into a set of relations.

FIGURE 9.8
An ER schema for a SHIP_TRACKING database.
Exercise 9.9 : Map this schema into a set of relations

FIGURE 9.9
EER diagram for a car dealer
Map conceptual schema design in the ER model to a relational database schema
  - Algorithm for ER-to-relational mapping
  - Illustrated by examples from the COMPANY database

Include additional steps in the algorithm for mapping constructs from EER model into relational model
Summary (2/2) – Mapping Steps

- **ER-to-Relational Mapping Algorithm**
  - Step 1: Mapping of Regular Entity Types
  - Step 2: Mapping of Weak Entity Types
  - Step 3: Mapping of Binary 1:1 Relation Types
  - Step 4: Mapping of Binary 1:N Relationship Types.
  - Step 5: Mapping of Binary M:N Relationship Types.
  - Step 6: Mapping of Multivalued attributes.
  - Step 7: Mapping of N-ary Relationship Types.

- **Mapping EER Model Constructs to Relations**
  - Step 8: Options for Mapping Specialization or Generalization.
  - Step 9: Mapping of Union Types (Categories).
Agenda

- The Role of Information Systems in Organizations
- The Database Design and Implementation Process
- Use of UML Diagrams as an Aid to Database Design Specification
- Rational Rose: A UML-Based Design Tool
- Automated Database Design Tools
Design methodology
  - Target database managed by some type of database management system

Various design methodologies

Large database
  - Several dozen gigabytes of data and a schema with more than 30 or 40 distinct entity types
Organizational context for using database systems

- Organizations have created the position of database administrator (DBA) and database administration departments
- Information technology (IT) and information resource management (IRM) departments
  - Key to successful business management
Database systems are integral components in computer-based information systems.

Personal computers and database system-like software products:
- Utilized by users who previously belonged to the category of casual and occasional database users.

**Personal databases** gaining popularity.

Databases are distributed over multiple computer systems:
- Better local control and faster local processing.
- **Data dictionary systems or information repositories**
  - Mini DBMSs
  - Manage *meta-data*

- **High-performance transaction processing systems require around-the-clock nonstop operation**
  - Performance is critical
Information system (IS)

Resources involved in collection, management, use, and dissemination of information resources of organization
- **Macro life cycle**
  - Feasibility analysis
  - Requirements collection and analysis
  - Design
  - Implementation
  - Validation and acceptance testing
  - Requirements collection and analysis
The database application system life cycle: **micro life cycle**

- System definition
- Database design
- Database implementation
- Loading or data conversion
- Application conversion
- Testing and validation
- Operation
- Monitoring and maintenance
- Design logical and physical structure of one or more databases
  - Accommodate the information needs of the users in an organization for a defined set of applications
- Goals of database design
  - Very hard to accomplish and measure
- Often begins with informal and incomplete requirements
Main phases of the overall database design and implementation process:

1. Requirements collection and analysis
2. Conceptual database design
3. Choice of a DBMS
4. Data model mapping (also called logical database design)
5. Physical database design
6. Database system implementation and tuning
Phases of Database Design and Implementation for Large Databases

Figure 10.1
Phases of database design and implementation for large databases.

**Phase 1:** Requirements collection and analysis

**Phase 2:** Conceptual database design

**Phase 3:** Choice of DBMS

**Phase 4:** Data model mapping (logical design)

**Phase 5:** Physical design

**Phase 6:** System implementation and tuning

- **Data content, structure, and constraints**
  - Data requirements
  - Conceptual Schema design (DBMS-independent)
  - Logical Schema and view design (DBMS-dependent)
  - Internal Schema design (DBMS-dependent)
  - DDL statements

- **Database applications**
  - Processing requirements
  - Transaction and application design (DBMS-independent)
  - Frequencies, performance constraints
  - Transaction and application implementation
Parallel activities

- **Data content, structure, and constraints** of the database
- Design of database applications

**Data-driven** versus **process-driven** design

**Feedback loops** among phases and within phases are common
Heart of the database design process

- Conceptual database design (Phase 2)
- Data model mapping (Phase 4)
- Physical database design (Phase 5)
- Database system implementation and tuning (Phase 6)
Activities

- Identify application areas and user groups
- Study and analyze documentation
- Study current operating environment
- Collect written responses from users
Requirements specification techniques

- Oriented analysis (OOA)
- Data flow diagrams (DFDs)
- Refinement of application goals
- Computer-aided
Phase 2a: Conceptual Schema Design

- Important to use a conceptual high-level data model
- Approaches to conceptual schema design
  - Centralized (or one shot) schema design approach
  - View integration approach
Strategies for schema design
- Top-down strategy
- Bottom-up strategy
- Inside-out strategy
- Mixed strategy

Schema (view) integration
- Identify correspondences/conflicts among schemas:
  - Naming conflicts, type conflicts, domain (value set) conflicts, conflicts among constraints
- Modify views to conform to one another
- Merge of views and restructure
Phase 2: Conceptual Database Design (3/3)

- Strategies for the view integration process
  - Binary ladder integration
  - N-ary integration
  - Binary balanced strategy
  - Mixed strategy

- Phase 2b: Transaction Design
  - In parallel with Phase 2a
  - Specify transactions at a conceptual level
  - Identify input/output and functional behavior
  - Notation for specifying processes
Phase 3: Choice of a DBMS

- Costs to consider
  - Software acquisition cost
  - Maintenance cost
  - Hardware acquisition cost
  - Database creation and conversion cost
  - Personnel cost
  - Training cost
  - Operating cost
- Consider DBMS portability among different types of hardware
Phase 4: Data Model Mapping (Logical Database Design)

- Create a conceptual schema and external schemas
  - In data model of selected DBMS
- Stages
  - System-independent mapping
  - Tailoring schemas to a specific DBMS
Phase 5: Physical Database Design

- Choose specific file storage structures and access paths for the database files
  - Achieve good performance

- Criteria used to guide choice of physical database design options:
  - Response time
  - Space utilization
  - Transaction throughput
Phase 6: Database System Implementation and Tuning

- Typically responsibility of the DBA
  - Compose DDL
  - Load database
  - Convert data from earlier systems
- Database programs implemented by application programmers
- Most systems include monitoring utility to collect performance statistics
Use of UML Diagrams as an Aid to Database Design Specification

- Use UML as a design specification standard
- Unified Modeling Language (UML) approach
  - Combines commonly accepted concepts from many object-oriented (O-O) methods and methodologies
  - Includes **use case diagrams**, **sequence diagrams**, and **statechart diagrams**
Advantages of UML

- Resulting models can be used to design relational, object-oriented, or object-relational databases
- Brings traditional database modelers, analysts, and designers together with software application developers
Different Types of Diagrams in UML (1/4)

- Structural diagrams
  - Class diagrams and package diagrams
  - Object diagrams
  - Component diagrams
  - Deployment diagrams
Different Types of Diagrams in UML (2/4)

- Behavioral diagrams
  - Use case diagrams
  - Sequence diagrams
  - Collaboration diagrams
  - Statechart diagrams
  - Activity diagrams
Use Case Diagram Notation

Figure 10.7
The use case diagram notation.
Different Types of Diagrams in UML (3/4)

Figure 10.9
The sequence diagram notation.
Different Types of Diagrams in UML (4/4)

**Figure 10.10**
The statechart diagram notation.

State consists of three parts:
- Name
- Activities
- Embedded machine
Activities and embedded machine are optional
Modeling and Design Example: UNIVERSITY Database

Figure 10.11
A sample statechart diagram for the UNIVERSITY database.
Figure 10.12
A sequence diagram for the UNIVERSITY database.
Figure 10.13
The design of the UNIVERSITY database as a class diagram.
Rational Rose: A UML-Based Design Tool

- **Rational Rose for database design**
  - Modeling tool used in the industry to develop information systems

- **Rational Rose data modeler**
  - Visual modeling tool for designing databases
  - Provides capability to:
    - **Forward engineer** a database
    - **Reverse engineer** an existing implemented database into conceptual design
- Reverse engineering
  - Allows the user to create a conceptual data model based on an existing database schema specified in a DDL file
- Forward engineering and DDL generation
  - Create a data model directly from scratch in Rose
  - Generate DDL for a specific DBMS
- Conceptual design in UML notation
  - Build ER diagrams using class diagrams in Rational Rose
- **Identifying relationships**
  - Object in a child class cannot exist without a corresponding parent object
- **Non-identifying relationships**
  - Specify a regular association (relationship) between two independent classes
Converting logical data model to object model and vice versa
- Logical data model can be converted to an object model
- Allows a deep understanding of relationships between conceptual and implementation models
- Synchronization between the conceptual design and the actual database
- Extensive domain support
  - Create a standard set of user-defined data types
- Easy communication among design teams
  - Application developer can access both the object and data models
Many CASE (computer-aided software engineering) tools for database design

Combination of the following facilities
- Diagramming
- Model mapping
- Design normalization
Characteristics that a good design tool should possess:

- Easy-to-use interface
- Analytical components
- Heuristic components
- Trade-off analysis
- Display of design results
- Design verification
Variety of products available

- Some use expert system technology

Table 10.1 Some of the Currently Available Automated Database Design Tools

<table>
<thead>
<tr>
<th>Company</th>
<th>Tool</th>
<th>Functionality</th>
</tr>
</thead>
<tbody>
<tr>
<td>Embarcadero Technologies</td>
<td>ER/Studio</td>
<td>Database modeling in ER and IDEF1x</td>
</tr>
<tr>
<td></td>
<td>DBArtisan</td>
<td>Database administration and space and security management</td>
</tr>
<tr>
<td>Oracle</td>
<td>Developer 2000 and</td>
<td>Database modeling, application development</td>
</tr>
<tr>
<td></td>
<td>Designer 2000</td>
<td></td>
</tr>
<tr>
<td>Persistence Inc.</td>
<td>PowerTier</td>
<td>Mapping from O-O to relational model</td>
</tr>
<tr>
<td>Platinum Technology (</td>
<td>Platinum ModelMart,</td>
<td>Data, process, and business component modeling</td>
</tr>
<tr>
<td>Computer Associates)</td>
<td>ERwin, BPwin, AllFusion Component</td>
<td></td>
</tr>
<tr>
<td>Popkin Software</td>
<td>Telelogic System Architect</td>
<td>Data modeling, object modeling, process modeling, structured analysis/design</td>
</tr>
<tr>
<td>Rational (IBM)</td>
<td>Rational Rose</td>
<td>Modeling in UML and application generation in C++ and Java</td>
</tr>
<tr>
<td>Resolution Ltd.</td>
<td>XCase</td>
<td>Conceptual modeling up to code maintenance</td>
</tr>
<tr>
<td>Sybase</td>
<td>Enterprise Application Suite</td>
<td>Data modeling, business logic modeling</td>
</tr>
<tr>
<td>Visio</td>
<td>Visio Enterprise</td>
<td>Data modeling, design and reengineering Visual Basic and Visual C++</td>
</tr>
</tbody>
</table>
Six phases of the design process
- Commonly include conceptual design, logical design (data model mapping), physical design

UML diagrams
- Aid specification of database models and design

Rational Rose and the Rose Data Modeler
- Provide support for the conceptual design and logical design phases of database design
A Case Study

- Implementing an ER diagram as a relational schema (relational database)
- General implementation of strong entities
- Handling attributes of different types
- General implementation of relationships
- Possible special implementation of binary many-to-one relationships
- Implementation of ISA
- Implementation of weak entities Foreign keys
- Primary key / foreign key constraints inducing many-to-one relationships between tables
- Concept of referential integrity
- Crow’s feet notation: ends of lines
- Crow’s feet notation: pattern of lines
We are now ready to convert ER diagrams into relational databases.

**Generally, but not always**
- An entity set is converted into a table
- A relationship is converted into a table

We will first go through a simple example.

Then, we will go through our large example, studied previously.

Then, we look at some additional points of interest.

Finally, we summarize the process, so we are sure we understand it.
Small ER Diagram

Employee
  ID#
  Name
  Child
  Likes
    Animal
      Species
      Discovered
  Born
    Country
      Name
      Population
More About The Example

- The given ER diagram is clear, other than
  - Discovered, which is the continent in which a particular species was first discovered
- Each child is a “dependent” of only one employee in our database
  - If both parents are employees, the child is “assigned” to one of them
- We are given additional information about the application
  - Values of attributes in a primary key must not be missing (this is a general rule, not only for this example)
  - Other than attributes in a primary key, other attributes unless stated otherwise may be missing
  - The value of Name is known for every Employee
- To build up our intuition, let’s look at some specific instance of our application
There are four countries, listing for them: Cname, Population (the latter only when known):

- US
- IN, 1150
- CN, 1330
- RU

We create a table for Country “in the most obvious way,” by *creating a column for each attribute* (underlying the attributes of the primary key) and this works:
There are five animals, listing for them: Species, Discovered (note, that even though not required, Discovered happens to be known for every Species):

- Horse, Asia
- Wolf, Asia
- Cat, Africa
- Yak, Asia
- Zebra, Africa

We create a table for Animal as before, and this works:
- There are five employees, listing for them: ID#, Name, (name of) Child (note there may be any number of Child values for an Employee, zero or more):
  - 1, Alice, Erica, Frank
  - 2, Bob, Bob, Frank
  - 4, Carol
  - 5, David
  - 6, Bob, Frank

- We create a table for Employee in the most obvious way, and this does not work:

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
<th>Child</th>
<th>Child</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Alice</td>
<td>Erica</td>
<td>Frank</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
<td>Bob</td>
<td>Frank</td>
</tr>
<tr>
<td>4</td>
<td>4</td>
<td>Carol</td>
<td></td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>5</td>
<td>David</td>
<td></td>
<td></td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Bob</td>
<td>Frank</td>
<td></td>
</tr>
</tbody>
</table>
- Child is a multivalued attribute so, the number of columns labeled “Child” is, in principle, unbounded.
- A table **must have** a fixed number of columns.
  - It must be an instance in/of a relational schema.
- If we are ready to store up to 25 children for an employee and create a table with 25 columns for children, perhaps tomorrow we get an employee with 26 children, who will not “fit”.
- We **replace our attempted single table** for Employee **by two tables**.
  - One for all the attributes of Employee other than the multivalued one (Child).
  - One for pairs of the form (primary key of Employee, Child).
- Note that both tables have a fixed number of columns, no matter how many children an employee has.
## Replace (incorrect)

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
<th>Child</th>
<th>Child</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Alice</td>
<td>Erica</td>
<td>Frank</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
<td>Bob</td>
<td>Frank</td>
</tr>
<tr>
<td>4</td>
<td>4</td>
<td>Carol</td>
<td></td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>5</td>
<td>David</td>
<td></td>
<td></td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Bob</td>
<td></td>
<td>Frank</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Alice</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
</tr>
<tr>
<td>4</td>
<td>4</td>
<td>Carol</td>
</tr>
<tr>
<td>5</td>
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<td>6</td>
<td>6</td>
<td>Bob</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Child</th>
<th>ID#</th>
<th>Child</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Erica</td>
</tr>
<tr>
<td>1</td>
<td>2</td>
<td>Frank</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Frank</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Frank</td>
</tr>
</tbody>
</table>
The primary key of the table Employee is ID#

The primary key of the table Child is the pair: ID#, Child

One attribute is not sufficient to get a primary key for Child

It is clear from the example how to handle any number of multivalued attributes an entity has

» Create a “main” table with all the attributes other than multivalued
   Its primary key is the original primary key of the entity set

» Create a table for each multivalued attribute consisting a primary key for the main table and that multivalued attribute
   Its primary key is the primary key of the entity combined with the multivalued attribute
Let us return to our example

Note that any value of ID# that appears in Child must also appear in Employee

» Because a child must be a dependant of an existing employee

This is an instance of a foreign key

ID# in Child is a foreign key referencing Employee

» This means that ID# appearing in Child must appear in some row “under” columns (here only one) of primary key in Employee

» Note that ID# is not a key of Child, so a foreign key in a table does not have to be a key of that table

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
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<tbody>
<tr>
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<td>6</td>
<td>Bob</td>
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</tbody>
</table>

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<thead>
<tr>
<th>Child</th>
<th>ID#</th>
<th>Child</th>
</tr>
</thead>
<tbody>
<tr>
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<tr>
<td>1</td>
<td>1</td>
<td>Frank</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Frank</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Frank</td>
</tr>
</tbody>
</table>
Foreign Key Induces A Many-To-One Relationship Between Tables

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
</tr>
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<tbody>
<tr>
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<tr>
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<td>1</td>
<td>Frank</td>
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</tr>
<tr>
<td></td>
<td>2</td>
<td>Frank</td>
</tr>
<tr>
<td></td>
<td>6</td>
<td>Frank</td>
</tr>
</tbody>
</table>

- **Note:**
  - Every row of Child has a single value of a primary key of Employee, so every row of Child “maps” to a single row of Employee
  - Every row of Employee has zero or more rows of Child mapped into it

In other words, no constraint
Likes needs to specify which employees like which animals

Such specification can be done using the primary keys of the entities

We do not need other attributes such as Name or Discovered

The table for likes contains some tuples:

- 1 likes Horse
- 1 likes Cat
- 2 likes Cat
- 6 likes Yak

<table>
<thead>
<tr>
<th>Likes</th>
<th>ID#</th>
<th>Species</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Horse</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>Cat</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Cat</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Yak</td>
</tr>
</tbody>
</table>
We could phrase the above somewhat differently

Likes needs to specify which employees (as rows in table Employee) like which animals (as rows in table Animal)

Such *a specification can done using the primary keys* of the tables Employee and Animal

The table for Likes contains some tuples:

- 1 likes Horse
- 1 likes Cat
- 2 likes Cat
- 6 likes Yak
Note that *there are foreign key constraints*

- ID# appearing in Likes is a foreign key referencing Employee
- Species appearing in Likes is a foreign key referencing Animal

And two many-to-one mappings are induced

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Alice</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
</tr>
<tr>
<td>4</td>
<td>4</td>
<td>Carol</td>
</tr>
<tr>
<td>5</td>
<td>5</td>
<td>David</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Bob</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Likes</th>
<th>ID#</th>
<th>Species</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Horse</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>Cat</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Cat</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Yak</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Animal</th>
<th>Species</th>
<th>Discovered</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Horse</td>
<td>Asia</td>
</tr>
<tr>
<td></td>
<td>Wolf</td>
<td>Asia</td>
</tr>
<tr>
<td></td>
<td>Cat</td>
<td>Africa</td>
</tr>
<tr>
<td></td>
<td>Yak</td>
<td>Asia</td>
</tr>
<tr>
<td></td>
<td>Zebra</td>
<td>Africa</td>
</tr>
</tbody>
</table>

This is true whenever we build a table for a relationship

- Likes was a relationship in the ER diagram
Born needs to specify which employees were born in which countries (for whom this information is known)

Such specification can done using the primary keys of the entities/tables

The relation Born contains some tuples:

- 1, US
- 2, IN
- 5, IN
- 6, CN

<table>
<thead>
<tr>
<th>Born</th>
<th>ID#</th>
<th>Cname</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>US</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>IN</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>IN</td>
<td></td>
</tr>
<tr>
<td>6</td>
<td>CN</td>
<td></td>
</tr>
</tbody>
</table>
Note that **there are foreign key constraints**

- ID# appearing in Born is a foreign key referencing Employee
- Cname appearing in Born is a foreign key referencing Country

And two many-to-one mappings are induced

- One of them happens to be one-to-one as an employee can be born in only one country
- This follows from the fact that in the ER diagram Born was a many-to-one relationship
- Compare with Likes, where an employee can like more than one animal and an animal can be liked by more than one Employee
- Let us focus on ID# in Employee
- No two different tuples in Born can have the same ID#
- Therefore ID# serves as a primary key, and we do not need Cname as part of the primary key

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
<th>Born</th>
<th>ID#</th>
<th>Cname</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Alice</td>
<td>1</td>
<td>US</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
<td>2</td>
<td>IN</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>5</td>
<td>Carol</td>
<td>5</td>
<td>IN</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>6</td>
<td>David</td>
<td>6</td>
<td>CN</td>
<td></td>
</tr>
<tr>
<td>6</td>
<td></td>
<td>Bob</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Country</th>
<th>CName</th>
<th>Population</th>
</tr>
</thead>
<tbody>
<tr>
<td>US</td>
<td></td>
<td></td>
</tr>
<tr>
<td>IN</td>
<td>1150</td>
<td></td>
</tr>
<tr>
<td>CN</td>
<td>1330</td>
<td></td>
</tr>
<tr>
<td>RU</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Using Visio

- Visio can be used to designing/specifying relational databases
- You can look at a tutorial, to get familiar with the mechanics of Visio
- This is greatly oversimplified, but a good start
  - [http://www.youtube.com/watch?v=1BYt3wmkgXE](http://www.youtube.com/watch?v=1BYt3wmkgXE) but foreign keys are not explained
  - [http://www.youtube.com/watch?v=55TpWp4TmMw&NR=1](http://www.youtube.com/watch?v=55TpWp4TmMw&NR=1)
  - [http://www.youtube.com/watch?v=r0x8ZMyPoj4&NR=1](http://www.youtube.com/watch?v=r0x8ZMyPoj4&NR=1) but this third part
    - Is misleading in the context of relational databases, due to the handling of many-to-many relationships and
    - The use of the second page, all the pages in a single Visio drawing refer to a *single* ER diagram, so each ER diagram needs its own Visio drawing/file
Specifying A Relational Implementation

- It is possible to use Visio or ErWin to specify our relational implementation
  - Visio has an “enterprise” version to generate database specifications from the diagram to SQL DDL
- We will just focus on the first task
- The second can be done automatically so we do not need to look at it here
Specifying A Relational Implementation (more on Visio)

- A drawing in Visio is not an Entity Relationship Diagram tool despite such terminology in Visio
  - In fact, this is good, as it produces a relational schema, which is what we actually need, but this is a lower-level construct than ER diagrams

- *It focuses on tables and the implicit many-to-one binary relationships induced by foreign key constraints*

- **Table**
  - A rectangle with three vertical subrectangles: name, list of attributes in the primary key, list of attributes not in the primary key
  - Required attributes are in bold
  - Attributes in the primary key and foreign keys are labeled as such

- **Relationship**
  - A many-to-one binary (or perhaps one-to-one, which is a special case) relationship induced by a foreign key constraint is explicitly drawn by means of a segment with an arrow head
  - We will have alternative notations later
## Relational Implementation For The Example

### Country

<table>
<thead>
<tr>
<th>Cname</th>
<th>Population</th>
<th>ID#</th>
<th>CName</th>
</tr>
</thead>
<tbody>
<tr>
<td>US</td>
<td></td>
<td>1</td>
<td>US</td>
</tr>
<tr>
<td>IN</td>
<td>1150</td>
<td>2</td>
<td>IN</td>
</tr>
<tr>
<td>CN</td>
<td>1330</td>
<td>5</td>
<td>IN</td>
</tr>
<tr>
<td>RU</td>
<td></td>
<td>6</td>
<td>CN</td>
</tr>
</tbody>
</table>

### Child

<table>
<thead>
<tr>
<th>ID#</th>
<th>Child</th>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Erica</td>
<td></td>
<td>1</td>
<td>Alice</td>
</tr>
<tr>
<td>1</td>
<td>Frank</td>
<td></td>
<td>2</td>
<td>Bob</td>
</tr>
<tr>
<td>2</td>
<td>Bob</td>
<td></td>
<td>4</td>
<td>Carol</td>
</tr>
<tr>
<td>2</td>
<td>Frank</td>
<td></td>
<td>5</td>
<td>David</td>
</tr>
<tr>
<td>6</td>
<td>Frank</td>
<td></td>
<td>6</td>
<td>Bob</td>
</tr>
</tbody>
</table>

### Animal

<table>
<thead>
<tr>
<th>Species</th>
<th>Discovered</th>
<th>Likes</th>
<th>ID#</th>
<th>Species</th>
</tr>
</thead>
<tbody>
<tr>
<td>Horse</td>
<td>Asia</td>
<td></td>
<td>1</td>
<td>Horse</td>
</tr>
<tr>
<td>Wolf</td>
<td>Asia</td>
<td></td>
<td>1</td>
<td>Cat</td>
</tr>
<tr>
<td>Cat</td>
<td>Africa</td>
<td></td>
<td>2</td>
<td>Cat</td>
</tr>
<tr>
<td>Yak</td>
<td>Asia</td>
<td></td>
<td>6</td>
<td>Yak</td>
</tr>
<tr>
<td>Zebra</td>
<td>Africa</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

### Diagram

- **Country**
  - PK: CName
  - FK: Country
  - CName: US
  - CName: IN
  - CName: CN
  - CName: RU
  - Population

- **Employee**
  - PK: ID#
  - FK: Child
  - ID#: Alice
  - ID#: Bob
  - ID#: Carol
  - ID#: David
  - ID#: Bob

- **Child**
  - PK,FK1: ID#
  - FK2: Child
  - ID#: 1
  - ID#: 2
  - ID#: 2
  - ID#: 6

- **Animal**
  - PK,FK1: ID#
  - FK2: Species
  - Species: Horse
  - Species: Cat
  - Species: Yak
  - Species: Zebra

- **Likes**
  - PK,FK1: ID#
  - FK2: Discovered
  - Discovered: Asia
  - Discovered: Africa
Cardinality Constraints

- The statement that a relationship is many-to-one as opposed to be a “standard” many-to-many relationship is really a cardinality constraint
- We will look at a relationships Likes between Person and Country and four cases of cardinality constraints on how many Countries a Person may like
  - No constraint
  - At least one
  - At most one
  - Exactly one
- For the first two, Likes is many-to-many
- For the last two, Likes is many-to-one
- Intuitively, Likes is many to one if for every Person, when you see which Countries this Person Likes, you get 0 or 1
- If you always get 1, this is a total function, otherwise this is a partial function
Specifying These Constraints (Revisited)

Every Person likes 0 or more Countries

Every Person likes 1 or more Countries

Every Person likes 0 or 1 Countries

Every Person likes 1 Country
Arrow Notation Cannot Distinguish Some Cases
Crow’s Feet: Improved Arrow Notation

- **Note:** different sides of the relationship are labeled in the two notations!
In general, cardinalities of both sides of the relationship may need to be specified.

We did only one, because it is sufficient to understand the notation.

We now return to the relational implementation of our example.

Visio and ErWin can use the Crow’s Feet notation.
## Relational Implementation For The Example

<table>
<thead>
<tr>
<th>Animal</th>
<th>Species</th>
<th>Discovered</th>
</tr>
</thead>
<tbody>
<tr>
<td>Horse</td>
<td>Asia</td>
<td>Asia</td>
</tr>
<tr>
<td>Wolf</td>
<td>Asia</td>
<td>Asia</td>
</tr>
<tr>
<td>Cat</td>
<td>Africa</td>
<td></td>
</tr>
<tr>
<td>Yak</td>
<td>Asia</td>
<td>Asia</td>
</tr>
<tr>
<td>Zebra</td>
<td>Africa</td>
<td></td>
</tr>
<tr>
<td>Zebra</td>
<td>Africa</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Born</th>
<th>ID#</th>
<th>CName</th>
</tr>
</thead>
<tbody>
<tr>
<td>Horse</td>
<td>1</td>
<td>US</td>
</tr>
<tr>
<td>Wolf</td>
<td>2</td>
<td>IN</td>
</tr>
<tr>
<td>Cat</td>
<td>5</td>
<td>IN</td>
</tr>
<tr>
<td>Yak</td>
<td>6</td>
<td>CN</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Child</th>
<th>ID#</th>
<th>Child</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Erica</td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>Frank</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>Bob</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>Frank</td>
<td></td>
</tr>
<tr>
<td>6</td>
<td>Frank</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Alice</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>Bob</td>
<td></td>
</tr>
<tr>
<td>4</td>
<td>Carol</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>David</td>
<td></td>
</tr>
<tr>
<td>6</td>
<td>Bob</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Likes</th>
<th>ID#</th>
<th>Species</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Horse</td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>Cat</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>Cat</td>
<td></td>
</tr>
<tr>
<td>6</td>
<td>Yak</td>
<td></td>
</tr>
</tbody>
</table>
Ends Of Lines (1/2)

- For every Country: 0..* of Born
  - 0 or more Employees were born there

- For every Born: 1..1 of Country
  - Because each row in Born has exactly 1 value of Country

- For every Employee: 0..1 of Born
  - Because an Employee was born in at most 1 Country

- For every Born: 1..1 of Employee
  - Because each row in Born has exactly 1 value of Employee

- For every Employee: 0..* Child
  - 0 or more of Child for an Employee

- For every Child: 1..1 Employee
  - Because every Child is assigned to exactly one Employee
For every Employee: 0..* of Likes
  » Employee can Like 0 or more Species

For every Likes: 1..1 of Employee
  » Because each row in Likes has exactly 1 value of Employee

For every Animal: 0..1 of Likes
  » Because a Species can be Liked by 0 or more Employees

For every Likes: 1..1 of Species
  » Because each row in Likes has exactly 1 value of Species
Born Versus Likes (1/2)

Country | Cname | Population
---------|-------|-------------
US       |       |             
IN       | 1150  |             
CN       | 1330  |             
RU       |       |             

Born | ID# | CName
-----|-----|-----
1    | US  |     
2    | IN  |     
5    | IN  |     
6    | CN  |     

Child | ID# | Child
------|-----|-----
1     | Erica
1     | Frank
2     | Bob
2     | Frank
6     | Frank

Employee | ID# | Name
--------|-----|-----
1       | Alice
2       | Bob
4       | Carol
5       | David
6       | Bob

Animal | Species | Discovered
-------|---------|-------------
Horse  | Asia    | Asia
Wolf   | Asia    | Asia
Cat    | Africa  | Africa
Yak    | Asia    | Asia
Zebra  | Africa  | Africa

Likes | ID# | Species
-----|-----|-----
1    | Horse
1    | Cat
2    | Cat
6    | Yak
Note that the many-to-one relationships are not of the same type in both cases.

The relationship between Likes and Employee indicates that when you start from a row of Employee you end up in between 0 and unbounded number of rows of Likes: no restriction.

An employee can like any number of animals.

The relationship between Born and Employee indicates that when you start from a row of Employee you end up in between 0 and 1 rows of Born.

An employee can be born in at most one country and therefore from a row of Employee you end up in between 0 and 1 rows of Born: a restriction.

*Born is really a (partial) one-to-one relationship*

Such relationships are considered “strange”
Treating Born Differently From Likes

- The above discussion implies that for every row in Employee there is at most one “relevant” row of Born
- Therefore, the “extra” information about an employee that is currently stored in Born can be added to Employee
- Born can be removed from the design
- This sounds very formal, but intuitively very clear as we can see from an alternative design
### Replace

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td></td>
<td>Alice</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>Bob</td>
</tr>
<tr>
<td>4</td>
<td></td>
<td>Carol</td>
</tr>
<tr>
<td>5</td>
<td></td>
<td>David</td>
</tr>
<tr>
<td>6</td>
<td></td>
<td>Bob</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Born</th>
<th>ID#</th>
<th>Cname</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td></td>
<td>US</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>IN</td>
</tr>
<tr>
<td>5</td>
<td></td>
<td>IN</td>
</tr>
<tr>
<td>6</td>
<td></td>
<td>CN</td>
</tr>
</tbody>
</table>

### by

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
<th>Cname</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td></td>
<td>Alice</td>
<td>US</td>
</tr>
<tr>
<td>2</td>
<td></td>
<td>Bob</td>
<td>IN</td>
</tr>
<tr>
<td>4</td>
<td></td>
<td>Carol</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td></td>
<td>David</td>
<td>IN</td>
</tr>
<tr>
<td>6</td>
<td></td>
<td>Bob</td>
<td>CN</td>
</tr>
</tbody>
</table>
Alternative Relational Implementation For The Example

<table>
<thead>
<tr>
<th>Country</th>
<th>CName</th>
<th>Population</th>
</tr>
</thead>
<tbody>
<tr>
<td>US</td>
<td></td>
<td></td>
</tr>
<tr>
<td>IN</td>
<td></td>
<td>1150</td>
</tr>
<tr>
<td>CN</td>
<td></td>
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</tr>
<tr>
<td>RU</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Child</th>
<th>ID#</th>
<th>Child</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Erica</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>Frank</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
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<td>2</td>
<td>Frank</td>
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<tr>
<td>6</td>
<td>6</td>
<td>Frank</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
<th>CName</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Alice</td>
<td>US</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
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</tr>
<tr>
<td>4</td>
<td>4</td>
<td>Carol</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>5</td>
<td>David</td>
<td>IN</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Bob</td>
<td>CN</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Animal</th>
<th>Species</th>
<th>Discovered</th>
</tr>
</thead>
<tbody>
<tr>
<td>Horse</td>
<td>Asia</td>
<td></td>
</tr>
<tr>
<td>Wolf</td>
<td>Asia</td>
<td></td>
</tr>
<tr>
<td>Cat</td>
<td>Africa</td>
<td></td>
</tr>
<tr>
<td>Yak</td>
<td>Asia</td>
<td></td>
</tr>
<tr>
<td>Zebra</td>
<td>Africa</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Likes</th>
<th>ID#</th>
<th>Species</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Horse</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>Cat</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Cat</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Yak</td>
</tr>
</tbody>
</table>
Alternative Relational Implementation For The Example

<table>
<thead>
<tr>
<th>Country</th>
<th>CName</th>
<th>Population</th>
</tr>
</thead>
<tbody>
<tr>
<td>US</td>
<td></td>
<td></td>
</tr>
<tr>
<td>IN</td>
<td></td>
<td>1150</td>
</tr>
<tr>
<td>CN</td>
<td></td>
<td>1330</td>
</tr>
<tr>
<td>RU</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Child</th>
<th>ID#</th>
<th>Child</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Erica</td>
</tr>
<tr>
<td>1</td>
<td>1</td>
<td>Frank</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Frank</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Frank</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Employee</th>
<th>ID#</th>
<th>Name</th>
<th>CName</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Alice</td>
<td>US</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Bob</td>
<td>IN</td>
</tr>
<tr>
<td>4</td>
<td>4</td>
<td>Carol</td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>5</td>
<td>David</td>
<td>IN</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Bob</td>
<td>CN</td>
</tr>
</tbody>
</table>

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<th>Species</th>
<th>Discovered</th>
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<td>Wolf</td>
<td>Asia</td>
<td></td>
</tr>
<tr>
<td>Cat</td>
<td>Africa</td>
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<td>Asia</td>
<td></td>
</tr>
<tr>
<td>Zebra</td>
<td>Africa</td>
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</tr>
</tbody>
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<th>ID#</th>
<th>Species</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>1</td>
<td>Horse</td>
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<tr>
<td>1</td>
<td>1</td>
<td>Cat</td>
</tr>
<tr>
<td>2</td>
<td>2</td>
<td>Cat</td>
</tr>
<tr>
<td>6</td>
<td>6</td>
<td>Yak</td>
</tr>
</tbody>
</table>
- The line between Animal and Likes is **solid** because the primary key of the “many side”, Likes, includes the primary key of the “one side”, Animal, so it “cannot exist” without it.
- The line between Employee and Likes is **solid** because the primary key of the “many side”, Likes, includes the primary key of the “one side”, Employee, so it “cannot exist” without it.
- The line between Employee and Child is **solid** because the primary key of the “many side”, Child, includes the primary key of the “one side”, Employee, so it “cannot exist” without it.
- The line between Country and Employee is **dashed** because the primary key of the “many side”, Employee, does not include the primary key of the “one side”, Country, so it “can exist” without it.
This is not a question of the ends of lines “forcing” the pattern of lines.

In the next slide, we see a slight modification of our example in which all lines have the same pair of endings.

We required that for each Employee the Country of Birth is known.

Nevertheless, as Cname is not part of the primary key of Country, the line is dashed.

For technical reasons, the tables have slightly different names, but this has nothing to do with our point.
Employee
PK ID#
Name
FK1 CName
Child
PK,FK1 ID#
PK Child
Country
PK CName
Population
Child1
PK,FK1 ID# Child
Employee1
PK ID#
FK1 Name CName
Animal
PK Species
Discovered
Likes
PK,FK1 ID# Species
Which Implementation To Use For Born?

- We cannot give a general rule
- The first implementation uses more tables
- The second implementation may introduce NULLs (empty values), which we do not like
- For the purpose of the class we will always use the second version, to have better exercises
- So do this for all the homeworks and tests, when relevant
Structurally, a relational database consists of:

1. A set of tables
2. A set of many-to-one binary relationships between them, induced by foreign key constraints.

In other words; a set of functions (in general partial), each from a table into a table.

When designing a relational database, you must specify both (or you will produce a bad specification).

- Technically, tables are enough, but this a very bad practice as you do not specify the relationships between tables.
- Tables are listed with attributes, specifying only which are in the primary key
  - Even the primary keys are not strictly required
- Foreign key constraints are not specified
  - So the DB system does not know what to enforce
- Even primary keys are not specified
We now convert our big ER diagram into a relational database

- We specify
  - Attributes that must not be NULL
  - Primary keys
  - Keys (beyond primary)
  - Foreign keys and what they reference
  - Cardinality constraints
  - Some additional “stubs”

- We both give a narrative description, similar to actual SQL DDL (so we are learning about actual relational databases) and Visio/Erwin diagrams

- We should specify domains also, but we would not learn anything from this here, so we do not do it

- *We go bottom up, in the same order as the one we used in constructing the ER diagram*
Hierarchy For Our ER Diagram

Note: circular dependency, need to be treated together

Note: circular dependency, need to be treated together

Note: circular dependency, need to be treated together

Type
Has
Loves
Student ISA
Professor ISA
Profession ISA
Monitors
Took
Taught
Section Offered
Prereq
Required
Car
Automobile
Person
Course
Book
Horse
Figure
Note
- Define Table Horse (Name NOT NULL, Primary Key (Name));
- This represents the simplest possible relational database
  - One table with one attribute
<table>
<thead>
<tr>
<th>PK</th>
<th>Name</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td></td>
</tr>
</tbody>
</table>
- Person has some interesting attributes
- Multivalued attribute: we will create another table
- Derived attribute: we do not create a column for it, it will be computed as needed
- Composite attribute: we “flatten” it
Define Table Person (ID# NOT NULL, SS# NOT NULL, FN, LN NOT NULL, DOB NOT NULL, Primary Key (ID#), Candidate Key (SS#), Age (computed by procedure …) );

- In SQL DDL, the keyword UNIQUE is used instead of Candidate Key, but “Candidate Key” is better for reminding us what this could be
- Age would likely not be stored but defined in some view
<table>
<thead>
<tr>
<th>PK</th>
<th>ID#</th>
</tr>
</thead>
<tbody>
<tr>
<td>SS#</td>
<td>FN</td>
</tr>
<tr>
<td>LN</td>
<td>DOB</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Horse</th>
</tr>
</thead>
<tbody>
<tr>
<td>PK</td>
</tr>
<tr>
<td>-------</td>
</tr>
</tbody>
</table>
Define Table Child (ID# NOT NULL, ChildName NOT NULL, Primary Key (ID#, ChildName), Foreign Key (ID#) References Person);

This lists all pairs (ID# of person, a child’s name)
  » We have chosen a more descriptive attribute name than the one in the ER diagram for children’s names

Note
  » A person may have several children, each with a different name
  » Two different persons may have children with the same name

Because of this, no single attribute can serve as primary key of Child
• Note that some attributes are not bold, such as FN here.
• This means that FN could be NULL (in this context, meaning empty).
• Note the induced many-to-one relationship.
• We need to make sure we understand what the line ends indicate:
  » A person may have 0 or more children (unbounded).
  » A child has exactly 1 person to whom it is attached.
• We need to pay attention to such matters, though we are generally not going to be listing them here.
But you should look at all lines and understand the ends and the patterns (solid or dashed).
Define Table Automobile (Model NOT NULL, Year NOT NULL, Weight NOT NULL, Primary Key (Model, Year));
Define Table Likes (ID# NOT NULL, Model NOT NULL, Year NOT NULL, Primary Key (ID#, Model, Year), Foreign Key (ID#) References Person, Foreign Key (Model, Year) References Automobile);
Define Table Car (VIN NOT NULL, Color, Primary Key (VIN));
There is no need for a table for Type as Type is a binary many-to-one relationship. It is essentially “stored” in the “many” side, that is in Car.
Define Table Car (VIN NOT NULL, Color, Model NOT NULL, Year NOT NULL, Weight NOT NULL, Primary Key (VIN), Foreign Key (Model, Year) References Automobile );
As Has is a binary many-to-one relationship, the attributed of this relationship, Date, is stored in the “many” side, Car

There is no need for a table for Has as Has is a binary many-to-one relationship

It is essentially “stored” in the “many” side, that is in Car

We can only specify that a Person has at least 1 Car with the notation we currently use

The CHECK condition is specified using appropriate SQL constraint syntax

This can actually be done in Visio/Erwin also
Define Table Car (VIN NOT NULL, Color, Model NOT NULL, Year NOT NULL, Weight NOT NULL, ID#, Primary Key (VIN), Foreign Key (Model, Year) References Automobile Foreign Key (ID#) References Person);
- We do not define a table for ISA
- This/these relationship/s is/are “embedded” in Student and Professor
Define Table Student (ID# NOT NULL, Primary Key (ID#), Foreign Key (ID#) References Person, GPA (computed by procedure …) );

Note, how ISA, the class/subclass (set/subset) relations, is modeled by Visio/Erwin
Define Table Professor ( 
ID# NOT NULL, 
Salary NOT NULL, 
Primary Key (ID#), 
Foreign Key (ID#) References Person );
Define Table Course (C# NOT NULL, Title NOT NULL, Description, Primary Key (C#));
Define Table Prereq (First NOT NULL, Second NOT NULL, Primary Key (First,Second), Foreign Key (First) References Course, Foreign Key (Second) References Course);
This is our first example of a table modeling a recursive relationship, between an entity set and itself.

We decide to name the table Prereq, as this is shorter than Prerequisite.

Note that it is perfectly clear and acceptable to refer here to C# by new names: First and Second.

» Similarly, to using ChildName in the Child table.

We should add some constraint to indicate that this (directed graph) should be acyclic.

» Maybe other conditions, based on numbering conventions specifying course levels.
Prerequisite (3/3)
Define Table Book (Author NOT NULL, Title NOT NULL, Primary Key (Author, Title));
Define Table Required (ID# NOT NULL, C# NOT NULL, Author NOT NULL, Title NOT NULL, Primary Key (ID#, C#, Author, Title), Foreign Key (ID#) References Professor, Foreign Key (C#) References Course, Foreign Key (Author, Title) References Book);

Why is it **bad** to have

Foreign Key (ID#) References Person, instead of

Foreign Key (ID#) References Professor?

Because only a Professor can Require a Book
This is our first example of a table modeling a relationship that is not binary.

Relationship Required was ternary: it involved three entity sets.

There is nothing unusual about handling it.

We still have as foreign keys the primary keys of the “participating” entities.
Define Table Section (C# NOT NULL, Year NOT NULL, Semester NOT NULL, Sec# NOT NULL, MaxSize, Primary Key (C#, Year, Semester, Sec#), Foreign Key (C#) References Course);

Note on the end of the edge between Course and Section, the Section end, on the drawing how the requirement of having at least one Section is modeled.
Section is our first example of a weak entity
- We do not define a table for Offered
- Relationship Offered is implicit in the foreign key constraint
- Define Table Took (ID# NOT NULL, C# NOT NULL, Year NOT NULL, Semester NOT NULL, Sec# NOT NULL, Grade,
  Primary Key (ID#, C#, Year, Semester, Sec#), Foreign Key (ID#) References Student, Foreign Key (C#, Year, Semester, Sec#) References Section);

- Note on the end of the edge between Section and Took, the Took end, on the drawing how the requirement of having between 3 and 50 students in a section is not fully modeled.

- We can only show 1 or more using current notation.
Because Took is a many-to-many relationship we store its attribute, Grade, in its table

We cannot store Grade in any of the two
» Section
» Student
Define Table Taught (ID# NOT NULL, C# NOT NULL, Year NOT NULL, Semester NOT NULL, Sec# NOT NULL, Primary Key (ID#, C#, Year, Semester, Sec#), Foreign Key (ID#), References Professor, Foreign Key (C#, Year, Semester, Sec#) References Section );
This is our first example in which a table, Taught, that “came from” a relationship is treated as if it came from an entity and participates in a relationship with other tables.

Nothing special needs to be done to “convert” a table that models a relationship, to be also treated as a table modeling an entity.

In this case, Monitors is a binary many-to-one relationship, so we do not need to create a table for it, and it can be stored in the “many” side, Taught.
Define Table Taught (ID# NOT NULL, C# NOT NULL, Year NOT NULL, Semester NOT NULL, Sec# NOT NULL, Monitor
Primary Key (ID#, C#, Year, Semester, Sec#),
Foreign Key (ID#), References Professor,
Foreign Key (C#, Year, Semester, Sec#) References Section
Foreign Key (Monitor) References Professor );
Additional Points

- We will discuss some additional, important, points
  - Elaboration on recursive relationships
  - Referential Integrity
  - Temporal databases
Assume now that a prerequisite course, “First” course, must be taken with at least some Grade to count as a prerequisite.

This to make an example a little “richer”

Two cases:

- A course may have any number of prerequisites
  Prereq is many-to-many
- A course may have at most one prerequisite
  Prereq is many to one (Second is the many side, a single First could be a prerequisite for many Second courses)
Nothing special, we handle the second case of Prereq by storing it in the “many” side of the relationship

So there are two additional attributes in Course1

» The prerequisite course, if any
» The required grade, if any
Assume that we have some professors in table Professor, with rows: 5,1 and 7,2

There is a row in Taught 5,G22.2433,2009,Spring,001,7

This means that 5 teaches a specific section and 7 monitors this assignment
Referential Integrity: Example (2/3)

- A user accesses the database and attempts to delete row (or all rows like this, recall that duplicates are permitted) 5,1 from **Professor**
- What should happen, as there is a row in Taught referencing this row in Professor?
- A user accesses the database and attempts to delete row 7,2 from **Professor**?
- What should happen, as there is a row in Taught referencing this row in Professor?
- Part of specification of foreign key in Taught
- An action on Professor can be denied, or can trigger an action on Taught
- For example
  - ON DELETE NO ACTION
    This means that the “needed” row in Professor cannot be deleted
    Of course, it is possible to delete the row from Taught and then from the Professor (if no other row in any table in the database “needs” the row in Professor)
  - ON DELETE CASCADE
    This means that if a row is deleted from Professor, all the rows in Taught referring to it are deleted too
  - ON DELETE SET NULL
    This means that the value referring to no-longer-existing professor is replaced by NULL
    In our example, this is not possible for ID# as it is a part of the primary key of Taught, but is possible for Monitor
Referential Integrity: Another Example

- Part of specification of foreign key in in Professor
- An action on Person can be denied, or can trigger an action on Professor
- For example
  - ON UPDATE CASCADE
    This means that if the value of ID# in Person is changed, this value of ID# also propagates to Professor
- Could (and probably should) add to Taught and Required:
  - ON UPDATE CASCADE
    In appropriate attributes, so that the change of ID# in Professor also propagates to them
    In Taught in both ID# and Monitor
    In Required in ID#
- Excellent mechanism for centralized maintenance
- Of course, we may want to maintain historical data
- So, in practice one may have some indication that the professor no longer works, but still keep historical information about the past
- But we do not assume this for our example
Example: **Person**

- Create a table for the entity without multivalued and derived attributes, flattening composite attributes.
- The primary key of this table will consist of the attributes serving as primary key of the entity.

Example table: Person

- If there is a derived attribute, describe how it is computed, but do not store it.
- If there is a multivalued attribute, create a table for it consisting of it and attributes of the primary key of the entity; do not put it in the table for the entity.

Example table: Child

The primary key of this table will consist of all its attributes.
There could be an attribute that is composite with some components being multivalued and some derived

And similar complexities

Example, without drawing the appropriate entity using the ER model (this is getting too hairy)

- A person has many children (multivalued)
- Each child has both FirstName and MiddleName
- The child has DOB
- The child has Age

Then the table for child will look like

<table>
<thead>
<tr>
<th>Child ID#</th>
<th>FirstName</th>
<th>MiddleName</th>
<th>DOB</th>
</tr>
</thead>
<tbody>
<tr>
<td>5432</td>
<td>Krishna</td>
<td>Satya</td>
<td>2006-11-05</td>
</tr>
</tbody>
</table>
Example: *ISA* and *Professor*

- Do not do anything for ISA
- The class “above” ISA (here Person) has already been implemented as a table
- Create a table with all the attributes of the subclass (as for strong entity above) augmented with the primary key of the table “above” ISA, and no other attributes from it

The primary key is the same as the primary key of the table “above” ISA

Example table: Professor
Example: *Offered* and *Section*

Do not do anything for the defining relationship, here Offered

Imagine that the weak entity is augmented by the primary key of the “stronger” table through which it is defined (the table for it has been created already)

Treat the augmented weak entity the same way as a strong entity

The primary key is the primary key of the “stronger” table augmented by the attributes in the discriminant of the weak entity (a discriminant may consist of more than one attribute)

Example table: Section and Offered
Example *Took*

The tables for the participating entities have already been created.

Create a table consisting of the primary keys of the participating tables and the attributes of the relationship itself.

Of course, treat attributes of the relationship that are derived, multivalued, or composite, appropriately, not storing them, producing additional tables, flattening them.

The primary key consists of all the attributes of the primary keys of the participating tables.

Example table: Took
Example: *Has*

Do not create a table for this relationship

Put the attributes of the primary key of the “one” side and the attributes of the relationship itself into the table of the “many” side

Of course, treat attributes of the relation that are derived, multivalued, or composite, appropriately, not storing them, producing additional tables, flattening them, as the case may be

You may decide to treat such a relationship the way you treat a relationship that is not binary many to one (but not in our class)

If the relationship is one-to-one, choose which side to treat as if it were “many”

Example table: Has
Example: *Taught* (before it was modified by removing Approved)

We have a table for that was created when we treated it as a relationship

We do not need to do anything else to this table

Example table: Taught
Agenda

1. Session Overview
2. ER and EER to Relational Mapping
3. Database Design Methodology and UML
4. Mapping Relational Design to ER/EER Case Study
5. Summary and Conclusion
Summary

- Basic ER model concepts of entities and their attributes
  - Different types of attributes
  - Structural constraints on relationships
- ER diagrams represent E-R schemas
- UML class diagrams relate to ER modeling concepts
- Enhanced ER or EER model
  - Extensions to ER model that improve its representational capabilities
  - Subclass and its superclass
  - Category or union type
- Notation and terminology of UML for representing specialization and generalization
Assignments & Readings

- **Readings**
  - Slides and Handouts posted on the course web site
  - Textbook: Chapters 9

- **Assignment #4**

- **Project Framework Setup (ongoing)**
Next Session: Relational Algebra, Relational Calculus, and SQL
Any Questions?